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Minimum degree and disjoint cycles in generalized claw-free graphs



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ABSTRACT

For $s \geq 3$ a graph is $K_{1,s}$ -free if it does not contain an induced subgraph isomorphic to $K_{1,s}$. Cycles in $K_{1,3}$ -free graphs, called claw-free graphs, have been well studied. In this paper we extend results on disjoint cycles in claw-free graphs satisfying certain minimum degree conditions to $K_{1,s}$ -free graphs, normally called generalized claw-free graphs. In particular, we prove that if G is $K_{1,s}$ -free of sufficiently large order $n = 3k$ with $\delta(G) \geq n/2 + c$ for some constant $c = c(s)$, then G contains k disjoint triangles. Analogous results with the complete graph K_3 replaced by a complete graph K_m for $m \geq 3$ will be proved. Also, the existence of 2-factors for $K_{1,s}$ -free graphs with minimum degree conditions will be shown.

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1. Introduction

In this paper we consider only graphs without loops or multiple edges. We let $V(G)$ and $E(G)$ denote the sets of vertices and edges of G , respectively. The *order* of G , usually denoted by n , is $|V(G)|$ and the *size* of G is $|E(G)|$. For any vertex v in G , let $N(v)$ denote the set of vertices adjacent to v and $N[v] = N(v) \cup v$. The *degree* $d(v)$ of a vertex v is $|N(v)|$, and we let $\delta(G)$ and $\Delta(G)$ denote the minimum degree and maximum degree of a vertex in G , respectively. If $U \subset V(G)$, we will use $G[U]$ to denote the subgraph of G induced by the vertices in U and let $E(U_1, U_2)$ denote the set of edges with one end in U_1 and one end in U_2 .

Let G and H be graphs. We say that G is H -free if H is not an induced subgraph of G . In this paper, we are interested in determining the number of disjoint cycles possible in a $K_{1,s}$ -free graph which satisfies certain minimum degree conditions.

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Disjoint cycles in claw-free graphs have been studied in a variety of papers. For example Chen, Faudree, Gould, and Saito investigated the range of the number of cycles in a 2-factor of a 2-connected claw-free graph G of order n with minimum degree $(n - 2)/3$ in [1].

Theorem 1. *If G is a 2-connected claw-free graph with $\delta(G) \geq \frac{n-2}{3}$, then G contains a 2-factor with exactly k cycles for $1 \leq k \leq \frac{n-24}{3}$. Furthermore, this result is sharp in the sense that if we lower $\delta(G)$ we cannot obtain the full range of values for k .*

Chen, Markus and Schelp studied independent cycles on the basis of edge density [2].

Theorem 2. *Let $k \geq 1$ and G be a $K_{1,s}$ -free graph of order n and size q .*

- (1) *If $s = 3$ and $q \geq \frac{1}{2}(3k - 1)(3k - 4) + 1$, then G contains k vertex disjoint cycles.*
- (2) *If $s \geq 4$ and $q \geq n16sk^2$, then G contains k disjoint cycles.*

The objective of this paper is to generalize the results for claw-free graphs proved in [3] to $K_{1,s}$ -free graphs for $s \geq 4$, and in particular to give analogues for the following three results.

Theorem 3. *Let k be a positive integer. If G is a claw-free graph of order*

$$n \geq 2k^4 - 2k^2 + k$$

with $\delta(G) \geq n/k$, then G contains a 2-factor with $k - 1$ components. Further, this value of $\delta(G)$ is best possible.

Theorem 4. *If G is a claw-free graph of order n with $\delta(G) \geq n/3$, then G contains a 2-factor with k disjoint cycles, for $2 \leq k \leq \lfloor n/3 - 2 \rfloor$.*

Theorem 5. *If G is a claw-free graph of sufficiently large order $n = 3k$ with $\delta(G) \geq n/2$, then G contains k disjoint triangles.*

We will need the following results in the proof of the main theorems. The next result, of Komlos, Sarkozy, and Szemerédi [4], verifies a conjecture of Seymour. A consequence of this result is that if G is a graph of sufficiently large order $n = r(k + 1)$ with $\delta(G) \geq kn/(k + 1)$, then G contains r vertex disjoint copies of K_{k+1} .

Theorem 6. *If $k \geq 1$ and G is a graph of sufficiently large order n with $\delta(G) \geq kn/(k + 1)$, then G contains the k th power of a Hamiltonian cycle.*

Ramsey numbers will be used in expressing the bounds on the number of vertex disjoint cycles and vertex disjoint complete graphs in a $K_{1,s}$ -free graph with varied minimum degrees. We will denote the Ramsey number $r(K_k, K_m)$ by the shorter notation $r(k, m)$.

Theorem 7 (Li, Rousseau and Zang [5]). *The Ramsey number*

$$r(K_k, K_n) \leq (1 + o(1)) \frac{n^{k-1}}{(\log n)^{k-2}}.$$

2. Disjoint complete graphs

The objective is to determine the number of possible disjoint complete graphs K_m for $m \geq 3$ in a $K_{1,s}$ -free graph with minimum degree at least n/k for some $k \geq 2$. The graph of Fig. 1 consists of k copies of the graph $K_{n/k}$ with an edge between two copies forming them into a ring. This graph has minimum degree $n/k - 1$. If $n/k = (t + 1)m - 1$, then $n = ktm + k(m - 1)$, but this graph will contain at most kt disjoint copies of a K_m . However, the order of the graph will accommodate as many as $kt + \lfloor \frac{k(m-1)}{m} \rfloor$ disjoint copies of a K_m . This implies that if G is a $K_{1,s}$ -free graph of order n and minimum degree at least n/k , then the maximum number of vertex disjoint copies of a K_m in G that will always exist will be at most $n/m - c$ for some constant $c = c(s, k)$. It will be shown that this does, in fact, always occur.

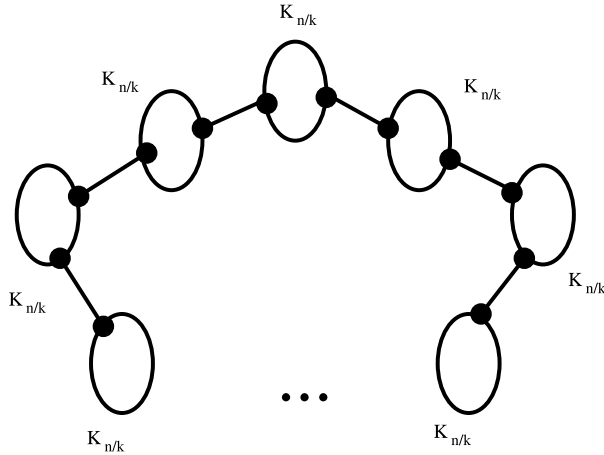


Fig. 1. $K_{1,s}$ -free graph G_1 of order $n = ktm + k(m - 1)$, $\delta \geq n/k$, but only kt disjoint K_m .

We begin with a look at disjoint triangles.

Theorem 8. For $s \geq 4$ and $r = r(3, s)$, let G be a $K_{1,s}$ -free graph of order n . If G has minimum degree δ , then G contains at least $F_3(n) = \left(\frac{3(\delta-s+1)}{3\delta+r-s-2}\right) \frac{n}{3}$ disjoint triangles.

Proof. Select a disjoint cycle system T composed of the maximum number, say t , of triangles. Let $H = G - V(T)$ be the subgraph of G that remains after removing T . No vertex of H can have degree s relative to H , since H is $K_{1,s}$ -free and contains no triangles. Thus for each $h \in V(H)$, $d_T(h) \geq \delta - s + 1$.

Consider a triangle $L \in T$ with vertices $\{x, y, z\}$ and let $\{a, b, c\}$ be the degrees of these vertices with respect to H respectively. We can assume with no loss of generality that $a \geq b \geq c$. We will show that $a + b + c \leq r + 2s - 5$. Assume not. If $a \geq r$, then $|N_H(x)| \geq r$, and since G is $K_{1,s}$ -free, there is a triangle in H , a contradiction. If $a < r$, then $b \geq s - 1$. Since G is $K_{1,s}$ -free, there is an edge in the neighborhood $N_{H \cup \{z\}}(y)$, and so there is a triangle L_1 with vertices y and two vertices of $N_{H \cup \{z\}}(y)$. Since $\lceil (r + 2s - 4)/3 \rceil \geq s + 1$, there is an edge in the neighborhood $N_{H \cup \{z\}}(x)$ that is disjoint from the vertices in L_1 . This implies that there is a triangle L_2 with vertices x and two vertices of $N_{H \cup \{z\}}(z)$ that are disjoint from L_1 . This contradicts the maximality of T . Thus, we can conclude that the vertices of each triangle in T collectively have at most $r + 2s - 5$ adjacencies in H .

The previous observation implies that $|E(T, H)| \leq t(r + 2s - 5)$, and so

$$(n - 3t)(\delta - s + 1) \leq |E(T, H)| \leq t(r + 2s - 5).$$

Thus,

$$(\delta - s + 1)n \leq (r - s - 2 + 3\delta)t;$$

hence,

$$t \geq \left(\frac{3(\delta - s + 1)}{3\delta + r - s - 2}\right) \frac{n}{3}. \quad \square$$

Consider the case when $\delta \geq n/k$ for $k \geq 2$. Thus,

$$t \geq \left(\frac{3(n/k - s + 1)}{3n/k + r - s - 2}\right) \frac{n}{3},$$

and so

$$t \geq \left(\frac{3n + k(r - s - 2)}{3n + k(r - s - 2)} - \frac{k(r + 2s - 5)}{3n + k(r - s - 2)}\right) \frac{n}{3}.$$

Therefore,

$$t \geq \frac{n}{3} - \left\lceil \frac{(r + 2s - 5)k}{9} \right\rceil.$$

Corollary 1. Let $s \geq 4, k \geq 2$, and $r = r(3, s)$. If G is a $K_{1,s}$ -free graph of sufficiently large order n with minimum degree $\delta(G) \geq n/k$ then G contains at least $\frac{n}{3} - \lceil \frac{(r+2s-5)k}{9} \rceil$ disjoint triangles.

Thus, for fixed s and k and n sufficiently large, a $K_{1,s}$ -free graph with minimum degree n/k has $n/3 - c$ vertex disjoint triangles for some constant $c = c(s, k)$. More specifically, if $s = 4$, then $r = r(3, 4) = 9$, and so we have the following bounds.

Corollary 2. If G is a $K_{1,4}$ -free graph of order n with minimum degree $\delta(G) \geq n/3$ then G contains at least $n/3 - 4$ disjoint triangles, and if the minimum degree $\delta(G) \geq n/2$ then G contains at least $n/3 - 3$ disjoint triangles.

In $K_{1,s}$ -free graphs, strong minimal degree conditions also imply the existence of many vertex disjoint copies of complete graphs K_m for $m \geq 4$. The following result, which is the analogue of Theorem 8, is an example of this.

Theorem 9. For $s \geq 4$ and $m \geq 4$ let G be a $K_{1,s}$ -free graph of order n . If G has minimum degree δ , then G contains at least $F_m(n) = \left(\frac{\delta - r(s, m - 1) + 1}{\delta - r(s, m - 1) + r(s, m)} \right) \frac{n}{m}$ disjoint copies of a complete graph K_m .

Proof. Select a disjoint system D composed of the maximum number, say d , of complete graphs K_m . Let $H = G - V(D)$ be the subgraph of G that remains after removing D . No vertex of H can have degree $r(s, m - 1)$ relative to H , since H is $K_{1,s}$ -free and does contain a copy of K_m . Thus for each $h \in H, d_D(h) \geq \delta - r(s, m - 1) + 1$.

If a vertex in D has as many as $r(s, m)$ adjacencies in H , then there would be a K_m in H , a contradiction. Thus, the number of edges between a $K_m \in D$ and H will be no more than $m(r(s, m) - 1)$.

The previous observations imply that

$$(n - dm)(\delta - r(s, m - 1) + 1) \leq |E(D, H)| \leq dm(r(s, m) - 1).$$

Thus,

$$(\delta - r(s, m - 1) + 1)n \leq dm((r(s, m) - 1) + \delta - r(s, m - 1) + 1);$$

hence,

$$d \geq \left(\frac{\delta - r(s, m - 1) + 1}{\delta - r(s, m - 1) + r(s, m)} \right) \frac{n}{m}. \quad \square$$

Consider the case when $\delta \geq n/k$ for $k \geq 2$. Then, in general from Theorem 9,

$$d \geq \left(\frac{\delta - r(s, m - 1) + 1}{\delta - r(s, m - 1) + r(s, m)} \right) \frac{n}{m},$$

and thus,

$$d \geq \left(\frac{\delta - r(s, m - 1) + r(s, m)}{\delta - r(s, m - 1) + r(s, n)} - \frac{r(s, m) - 1}{\delta - r(s, m - 1) + r(s, m)} \right) \frac{n}{m},$$

or equivalently

$$d \geq \frac{n}{m} - \left(\frac{r(s, m)}{\delta - r(s, m - 1) + r(s, m)} \right) \frac{n}{m}.$$

Therefore, when $\delta = n/k$,

$$d \geq \frac{n}{m} - \left(\frac{nr(s, m)k}{mn - kmr(s, m - 1) + kmr(s, m)} \right),$$

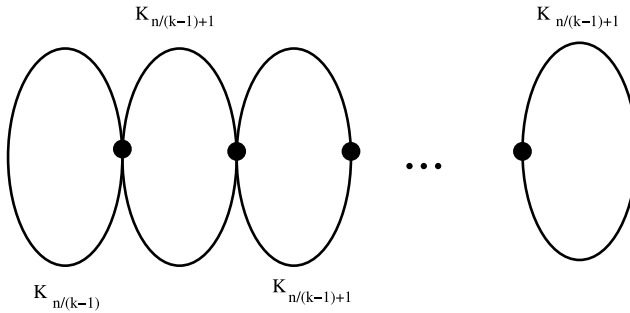


Fig. 2. G_2 composed of $k - 1$ blocks with no 2-factor with $k - 2$ cycles.

which implies

$$d \geq \frac{n}{m} - \left\lceil \left(\frac{r(s, m)k}{m} \right) \right\rceil,$$

since $r(s, m) - r(s, m - 1)$ is a positive integer. \square

This results in the following corollary.

Corollary 3. For $s \geq 4$ and $k \geq 2$ let G be a $K_{1,s}$ -free graph of order n . If G has minimum degree n/k , then G contains at least $\frac{n}{m} - c$ vertex disjoint copies of K_m for some $c = c(m, k, s)$. More specifically, G has at least $\frac{n}{m} - \lceil \frac{r(s,m)k}{m} \rceil$ vertex disjoint copies of K_m .

For example, a graph G of sufficiently large order n with minimum degree $n/4$ will have at least $n/4 - 18$ disjoint copies of a K_4 , since $r(4, 4) = 18$.

3. Disjoint cycles

The objective of this section is to determine the number of possible cycles in a 2-factor in a $K_{1,s}$ -free graph with minimum degree at least n/k for some $k \geq 2$. Consider the graph G_2 formed by taking one copy of $K_{n/(k-1)}$ and identifying a vertex with a vertex in a copy of $H_2 = K_{n/(k-1)+1}$. Now identify a new copy of H_2 with a different vertex of the last copy, and repeat this process until we have a “path” of subgraphs with $k - 1$ blocks (see Fig. 2). The graph G_2 is $K_{1,s}$ -free and has order n , and $\delta(G_2) = n/(k - 1) - 1$. Also, $n/(k - 1) - 1 \geq n/k$ whenever $n \geq (k - 1)k$, and G_2 clearly has a 2-factor with $k - 1$ components, but no 2-factor with $k - 2$ cycles.

To verify that a $K_{1,s}$ -free graph G of order n with $\delta(G) \geq n/k$ has a 2-factor with $k - 1$ components, we will need the following lemma on the independence number of such a graph.

Lemma 1. If G is a $K_{1,s}$ -free graph with $\delta(G) \geq n/k$ for $k \geq 2$, then the independence number $\alpha(G) \leq (s - 1)k - 1$.

Proof. Choose an independent set S with $\alpha = \alpha(G)$ vertices. Let $H = G - S$ be the remaining subgraph of order $n - \alpha$. Any vertex of H has degree at most $s - 1$ in S as G is $K_{1,s}$ -free. Further, each vertex of S has all its neighbors in H . If $E = E(S, H)$ is the set of edges between S and H , then

$$\alpha \binom{n}{k} \leq |E| \leq (s - 1)(n - \alpha)$$

and so

$$\alpha \leq \frac{(s - 1)kn}{n + (s - 1)k} = k(s - 1) \left(\frac{n}{n + (s - 1)k} \right) < (s - 1)k;$$

hence,

$$\alpha(G) \leq (s - 1)k - 1. \quad \square$$

Theorem 10. Let k be a positive integer, and $s \geq 4$. If G is a $K_{1,s}$ -free graph of sufficiently large order n with $\delta(G) \geq n/k$, then G contains a 2-factor with $k - 1$ components. Further, this value of $\delta(G)$ is best possible, in that $\delta(G) \geq n/(k + 1)$ is not sufficient.

Proof. Suppose we select a vertex disjoint set system \mathcal{C} with $k - 1$ cycles C_1, C_2, \dots, C_{k-1} , where $|\cup_{i=1}^{k-1} V(C_i)|$ is as large as possible. We know that such a set exists from [Corollary 1](#). Let $H = G - \cup_{i=1}^{k-1} V(C_i)$.

Observe that with any one cycle C_i , a vertex $h \in V(H)$ has at most $(s - 1)k - 1$ adjacencies, for otherwise there would exist an independent set (predecessors of adjacencies along with h) of order at least $(s - 1)k$, a contradiction to [Lemma 1](#). Thus, $\delta(H) \geq n/k - (k - 1)((s - 1)k - 1)$.

But the bound on $\delta(H)$ implies that H contains a cycle of length at least $\delta(H) + 1$. Thus, as \mathcal{C} is as large as possible, each cycle C_i ($1 \leq i \leq k - 1$) contains at least $\delta(H) + 1 \geq n/k - c'$ vertices for some constant $c' = c'(k, s)$. This, in turn, implies that $V(H) \leq n/k + c$ for some constant $c = c(s, k)$. Hence, for n sufficiently large, H is dense and, in fact, H is hamiltonian connected, since $2(n/k - c')$ is significantly larger than $n/k + c$.

Claim 1. No cycle in \mathcal{C} has two independent edges to H .

Suppose this were not the case; say, C_b has edges $w_i h_i$ and $w_j h_j$ with $w_i, w_j \in V(C_b)$ and $h_i, h_j \in V(H)$. Without loss of generality we can assume that w_i, w_{i+1}, \dots, w_j contains more than half of the vertices of C_b . Therefore, the cycle

$$(w_i, w_{i+1}, \dots, w_j, h_j, P, h_i, w_i),$$

where P is a hamiltonian path connecting h_i and h_j in H , is a cycle longer than C_b , contradicting our choice of \mathcal{C} .

Claim 2. No two cycles of \mathcal{C} have three independent edges between them.

Suppose instead that C_a and C_b had three independent edges between them. Without loss of generality say that $a_1 b_1, a_2 b_2$ and $a_3 b_3$ are these edges with $a_i \in C_a$ and $b_i \in C_b, i = 1, 2, 3$. Also, without loss of generality, suppose that the segment (a_1, a_2) contains less than $|C_a|/3$ vertices and (b_1, b_2) contains less than $|C_b|/2$ vertices. Then, a new cycle

$$C'_a = (a_2, a_2^+, \dots, a_1, b_1, b_1^-, \dots, b_2, a_2)$$

replaces C_a and H replaces C_b to form a new system with more vertices than \mathcal{C} , a contradiction.

By [Claims 1](#) and [2](#) we see that some cycles may have a vertex of large degree to H , but then no other vertices of that cycle have any adjacencies in H .

Observe that each vertex of H has edges to \mathcal{C} . If this were not true, and $d_{\mathcal{C}}(h) = 0$ for some $h \in V(H)$, then since $d(h) \geq n/k$, this implies that $|H| \geq n/k + 1$. Since every cycle in \mathcal{C} is at least as large as H , this gives the contradiction that $n \geq k(n/k + 1) = n + k$. By the same reasoning, no vertex of H has only one edge to \mathcal{C} , because if this were the case then we would have $|H| \geq n/k - 1 + 1 = n/k$ and, hence, $|H| = n/k = |C_i|$ for $i = 1, 2, \dots, k - 1$. But then every vertex of every cycle has edges to other cycles, which is in contradiction to one of the [claims 1](#) or [2](#).

The previous observations imply that each of the cycles C_i and H induce dense subgraphs of order approximately n/k . That is, with the exception of a function of $c^* = c^*(k, s)$ vertices in each cycle, the vertices have degree at least $n/k - c_1$ for some $c_1 = c_1(k, s)$. Since each cycle is only of order at most $n/k + c_2$ for some $c_2 = c_2(k, s)$, these dense subgraphs will have strong hamiltonian properties. For example, even after a small number of vertices are removed, a cycle will span the rest of the dense subgraph.

Now suppose that $H = \{C_0, C_1, \dots, C_q\}$ are the cycles with edges to other cycles. If we consider these cycles as the vertices of a graph, then among these $q + 1$ cycles there are at least $q + 1$ independent edges, and a cycle of cycles can be formed.

Say that $\{C_{i_1}, C_{i_2}, \dots, C_{i_t}, C_{i_1}\}$ are the “vertices” of this cycle. Then, starting in C_{i_1} we may traverse all but a function of k and s vertices before we cross to C_{i_2} . In C_{i_2} we traverse all but a function of k and s vertices before we cross to C_{i_3} , where we traverse a minimum number of vertices (some function k

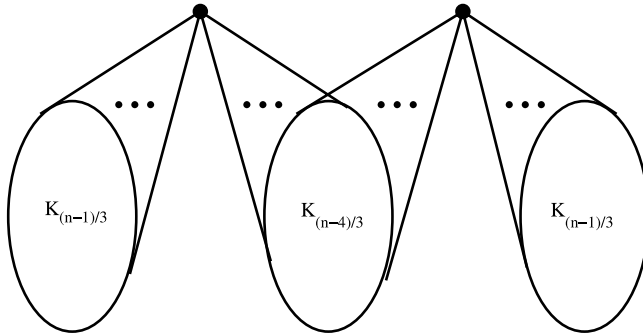


Fig. 3. Claw-free graph G_3 , with no 2-factor consisting of two cycles.

and s) before we cross to C_{i_4} . Continuing in this manner we return to C_{i_1} , completing a cycle. Now on the other subgraphs corresponding to this cycle we form new cycles using a maximum number of the remaining dense subgraphs. Thus, at most a function of k and s vertices has been lost from any of the original cycles.

We now form \mathcal{C}' to include all these new cycles, as well as H if it is not a part of these cycles, and all the unchanged cycles from \mathcal{C} . This is a system of $k - 1$ cycles that includes all but a function of k and s vertices of G , contradicting our choice of \mathcal{C} and completing the proof. \square

The graph G_2 in the case $k = 3$ shows that $\delta(G) \geq n/2$ is needed to obtain a Hamiltonian cycle in a $K_{1,s}$ -free graph of order n . The graph G_3 of Fig. 4 has order n and $\delta(G_3) = \frac{n-1}{3}$, but clearly cannot be covered by two cycles. Thus $\delta(G) \geq n/3$ is required to have a 2-factor with just two cycles (see Fig. 3).

Theorem 11. *If G is a $K_{1,s}$ -free graph of order n with $\delta(G) \geq n/3$, then G contains a 2-factor with k disjoint cycles for $2 \leq k \leq \lfloor n/3 - \frac{r(3,s)+2s-5}{3} \rfloor$.*

Proof. When $k = 2$, the result holds by Theorem 10. Suppose we select a disjoint cycle system $\mathcal{C} = \{C_1, C_2, \dots, C_t\}$ for each $t \geq 3$ in the range. We know that such a system exists by Corollary 1. Assume that \mathcal{C} is chosen to contain the maximum number of vertices, and let $H = G - \mathcal{C}$.

Observe that if $d_H(h) > n/(t + 1)$ for all $h \in V(H)$, then H contains a cycle of length greater than $n/(t + 1)$ and, hence, each cycle in \mathcal{C} has length greater than $n/(t + 1)$, or we could find a system larger than \mathcal{C} . This implies $|V(G)| = n > (t + 1)(n/(t + 1)) = n$, a contradiction. Therefore, for each $t \geq 3$ there exists a vertex $h \in V(H)$ such that $d_{\mathcal{C}}(h) \geq n/3 - n/(t + 1)$. We also have by Lemma 1 that $\alpha(G) \leq 3s - 4$.

Previous arguments imply that there is a vertex $x \in V(H)$ such that $d_{\mathcal{C}}(x) \geq cn$ for some constant c . Observe that x has at most $3s - 5$ adjacencies to any cycle of \mathcal{C} , since more adjacencies would imply an independent set with at least $3s - 3$ vertices using predecessors of the adjacencies of x and x . Therefore, x is adjacent to a function of n different cycles of \mathcal{C} , say q . Hence $q \geq cn/(3s - 5)$.

Let $X = \{x_1, x_2, \dots, x_q\}$ be the adjacencies of x in these q cycles. Since $\alpha(X) \leq 3s - 4$, there is a subset $X_1 \subset X$ that induces a complete graph and $|X_1| \geq q^{1/3s}$. Let X_1^+ be the predecessor of the vertices of X_1 on the respective cycles. There is a subset $X_2 \subset X_1^+$ that induces a complete graph. This can be repeated with the successors of the adjacencies of x to form a subset $X_3 \subset X_2$ with at least two vertices. This implies that there are vertices $y_1, y_2 \in X$ in cycles C' and C'' respectively such that $y_1 y_2 \in E(G)$, $y_1^+ y_2^+ \in E(G)$, and $y_1^- y_2^- \in E(G)$. The two cycles C' and C'' can be replaced by the cycle (x, y_1, y_2, x) and the cycle formed from $C' - \{y_1\}$ and $C'' - \{y_2\}$ using the edges $y_1^+ y_2^+$ and $y_1^- y_2^-$. This contradicts the maximality of the cycle system \mathcal{C} , and completes the proof of Theorem 4. \square

4. Complete graph factors

In [3] it was shown that in a claw-free graph of order $n = 3k$, $\delta(G) \geq n/2$ is sufficient to imply that there are k vertex disjoint triangles (Theorem 5). The minimum degree condition $\delta(G) \geq n/2$ is not sufficient if the triangle K_3 is replaced by the a complete graph K_m for $m \geq 4$ with n divisible by m .

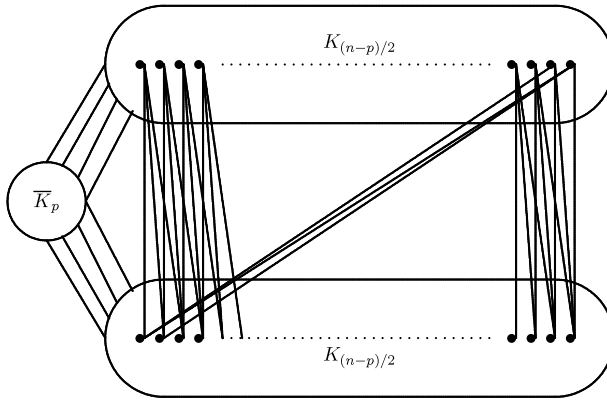


Fig. 4. G_4 .

For a fixed integer p with $n - p$ divisible by 2, consider the graph $\overline{K}_p + (K_{(n-p)/2} \cup K_{(n-p)/2})$. Let $X = \{x_1, x_2, \dots, x_{(n-p)/2}\}$ and $Y = \{y_1, y_2, \dots, y_{(n-p)/2}\}$ be the vertices of the two complete graphs. For $m \geq 4$ and for each i with $1 \leq i \leq (n - p)/2$ add the edges $x_i y_i, x_i y_{i+1}, \dots, x_i y_{m-3}$ with the indices taken modulo $(n - p)/2$. Denote this graph by G_4 (see Fig. 4). There is no K_m in G_4 with vertices in both X and Y , and so all copies of a K_m will have all of its vertices in either X or Y or $m - 1$ vertices in either X or Y and one vertex in \overline{K}_p . Therefore, if n is divisible by m , and there are n/m vertex disjoint copies of a K_m , then $p = p_1 + p_2$ such that $(n - p)/2 - p_i(m - 1)$ is divisible by m for $i = 1, 2$. This implies that $p(m - 2)$ is divisible by m . Hence, if p is chosen such that $p(m - 2)$ is not divisible by m and $p < s$, then G_4 does not contain n/m vertex disjoint copies of a K_m . However, $\delta(G_4) \geq (n + p - 8 + 2m)/2 > n/2$ for $m \geq 4$ and $p \geq 1$. Thus, a minimum degree condition of $\delta(G) \geq n/2 + c$ where $c = c(m, s)$ will be needed to imply the existence of n/m vertex disjoint copies of a K_m .

Our goal in this section is to prove the following result.

Theorem 12. *Let $m \geq 4$ and $s \geq 3$. If G is a $K_{1,s}$ -free graph of sufficiently large order $n = km$, then there is a $c = c(s, m)$ such that if $\delta(G) \geq n/2 + c$, G contains k disjoint copies of K_m .*

Proof of Theorem. By Lemma 1, $\alpha(G) \leq 2s - 3$. Since G does not contain $2s - 3$ independent vertices, Ramsey theory implies that G contains a large clique; in fact, G contains a $K_{\lfloor \frac{1}{n^{2s-2}} \rfloor}$. Select such a clique and denote it by A . Let $B \subseteq G - A$ be those vertices of $G - A$ whose degree to A is at most $r^* = m(r(m, 2s - 2) - 1)$. Let $C = G - (A \cup B)$.

Observe that

$$|E(A, C)| \geq |A|(n/2 + c - |A|) - r^*|B|.$$

Thus,

$$|C| \geq \frac{|A|(n/2 + c - |A|) - r^*|B|}{|A|},$$

since each vertex in C has at most A adjacencies in A . However, since $|A| \geq n^{\frac{1}{2s-2}}$, and c and r^* are constants and not a function of n ,

$$|C| \geq n/2 - o(n).$$

Let

$$B_2 = \{b \in B \mid d_C(b) \geq mr(m, 2s - 2)\},$$

and let $B_1 = B - B_2$. Note that each vertex in B_1 has at most $2(m - 1)r(m, 2s - 2)$ adjacencies in $A \cup C$ and so if B_1 is nonempty,

$$|B_1| \geq n/2 - 2mr(m, 2s - 2) \approx n/2.$$

Now we consider the partition $V(G) = B_1 \cup D$, where $D = A \cup B_2 \cup C$. Note that we have both $|B_1| \approx n/2$ and $|D| \approx n/2$. If $|B_1| \equiv 0 \pmod{m}$, then let $B'_1 = B_1$. If $|B_1| \geq n/2$, then every vertex of D must have at least c adjacencies to B_1 . Hence, as G is $K_{1,s}$ -free and $c = c(s, m)$ is large, we may find a K_m containing $m - 1$ vertices of B_1 and one vertex of D . Remove this copy of a K_m . Continue to do this until we get a subgraph B'_1 of B_1 such that

$$|B'_1| \equiv 0 \pmod{m}.$$

If $|B_1| < n/2$, then each vertex of B_1 has at least c adjacencies to D . As before, we can find a copy of K_m containing one vertex of B_1 and $m - 1$ vertices of D . Remove this K_m and continue this until we get a subgraph B'_1 of B_1 such that

$$|B'_1| \equiv 0 \pmod{m}.$$

Now, since B'_1 is very dense and has order a multiple of m , and n is sufficiently large, we may apply [Theorem 6](#) to B'_1 to obtain an independent set of disjoint copies of K_m that covers all of B'_1 .

We can find a copy of K_m in the vertices of B_2 as long as there are at least $r(m, 2s - 2)$ vertices remaining in B_2 . Each of the remaining vertices after the deletion of the K_m have at least $mr(m, 2s - 2)$ adjacencies in C , so each of these remaining vertices can be placed in a K_m using $m - 1$ vertices in C .

We can find a copy of K_m in the vertices of C as long as there are at least $r(m, 2s - 2)$ vertices remaining in C . Each of the remaining vertices after the deletion of the K_m have at least $(m - 1)r(m, 2s - 2)$ adjacencies in A , so each of these remaining vertices can be placed in a K_m using $m - 1$ vertices in A . Since A is a complete graph, the remaining vertices of A can be partitioned into disjoint copies of complete graphs K_m . \square

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